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Recursion(int day){return Recursion(day += 1);}

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Implementation

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"To create recursion, you must create recursion."

[Recursive](#page-1-0) design

[re-computing](#page-12-0) a loop

How to design a recursive algorithm

- **1** First write the base cases. Must always have some base cases, which can be solved without recursion.
- **2** Think about solving the problem by combining the results of one or more smaller, but similar, sub-problems. If the algorithm you write is correct, then certainly you can rely on it (recursively) to solve the smaller subproblems.
- **3** Making progress. For the cases that are to be solved recursively, the recursive call must always be to a case that makes a **fixed** quantity of progress toward a base case (not fixed proportion).
- 4 Compound interest guideline: Never duplicate work by solving the same instance of a problem in separate recursive calls.
- **6** Check progressively larger inputs to inductively validate that there is no infinite recursion
- **6** The secret to success is: Do not worry about how the recursive call solves the subproblem. Simply accept that it will solve it correctly, and use this result to in turn correctly solve the original problem.

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[Palindrome](#page-5-0) detection

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Example: Palindrome checking function

- Recursive case?
- Base case (which often results in termination)?
- Condition/test, which checks for the base?

Observe code to show recursion stack

CCD

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Non-recursive GCD: Euclid's algorithm

Euclid's algorithm efficiently computes the greatest common divisor (GCD) of two numbers (AB and CD below), the largest number that divides both without leaving a remainder (CF).

Proceeding left to right:

Recursive GCD

[GCD](#page-6-0)

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Check out the code

[Simple word](#page-8-0) reversal

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Very simple word reversal

[Simple word](#page-8-0) reversal

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Use recursion to use the cin/cout buffer to reverse inputted text: see code

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What is involved when making a function call?

Overhead:

- save caller's state
- allocate stack space for arguments
- allocate stack space for local variables
- invoke routine at exit (return), release resources
- restore caller's "environment"
- resume execution of caller

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Observe code

Some recursive functions can be unnecessarily inefficient, and would be better as iterative functions, e.g., recursive Fibonacci: $fib(n)$ is 0 if n is zero fib (n) is 1 if n is one fib(n) is fib(n - 1) + fib(n - 2) otherwise $\overbrace{}$ ß ß

Efficiency

Problem: [re-computing](#page-12-0) values

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Some recursive functions can be unnecessarily inefficient, and would be better as iterative functions, e.g., recursive Fibonacci: $fib(n)$ is 0 if n is zero $fib(n)$ is 1 if n is one

fib(n) is fib(n - 1) + fib(n - 2) otherwise

Observe code

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Use a loop

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If you can easily find a looping algorithm (e.g., fibonacci with a loop below:)

Remember, recursion is just another way to loop

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[Solution 2: Tail](#page-17-0) recursion calls

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[re-computing](#page-12-0) values a loop

[Tail recursion](#page-18-0)

- Iterative
-
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Tail recursion: just a glorified loop

- • Tail call is a subroutine call performed as the final action of a procedure (recall order mattering example)
- Tail calls don't have to add new stack frame to the call stack.
- Tail recursion is a special case of recursion where the calling function does no more computation after making a recursive call.
- Tail-recursive functions are functions in which all recursive calls are tail calls and thus do not build up any deferred operations.
- Producing such code instead of a standard call sequence is called tail call elimination.
- Tail call elimination allows procedure calls in tail position to be implemented as efficiently as goto statements, thus allowing more efficient structured programming.

values a loop

Iterative [conversion](#page-19-0)

stack

Convert an iterative loop into a tail recursive

- **1** First identify those variables that exist outside the loop but are changing in the loop body; these variable will become formal parameters in the recursive function.
- 2 One then builds a function that has these "outside" variables as formal parameters, with default initial values
- **3** The original loop test becomes an if() test in the body of the new function
- **4** The if-true block becomes the recursive call.
- **6** Arguments to the recursive call encode the updates to the loop variables.
- **6** else block becomes the value the loop attempted to calculate
- **2** Conversion results in tail recursion

Examples in code

Language choice

[Language](#page-20-0) choice

 gcc-c++ (aka g++) usually can do tail call optimization. Tail call elimination doesn't automatically happen in Java or Python, though it does reliably in functional languages like Lisp: Scheme, Clojure (upon specification), and Common Lisp, which primarily employ recursion.

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Solution 3: [Memoization /](#page-21-0) caching

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Memoization and caching

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Solution 3: [Memoization /](#page-21-0) caching

Store the values which have already been computed, rather than re-compute them:

See code for Fibonacci

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Convert recursion to loops

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{

} }

[Convert loops](#page-23-0) to recursion

Conversion of recursive to loops is less systematic than converting a loop to tail recursive, requires more creativity, and isn't always easy for a human, though converting an already tail recursive algorithm is more straightforward:

```
void tail (int i)\{if ( i > 0 ) {
  cout << i << '' ;
       tail (i - 1);
}
```
void iter Equivalent Of Tail (int i) { for $(; i > 0; i --)$ { cout $<<$ i $<<$ ' \lbrack ':

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Implementing recursion with a stack

[re-computing](#page-12-0)

Implementing [recursion with a](#page-25-0) stack

Often converting more inherently recursive algorithms to loops requires keeping a programmer-designed stack like would be done by the compiler in the first place. See example for factorial.